


```
1  xbegin
2   read a
3   xbegin_open
4     read i
5     i ← i + 1
A 6   write i
7   xend_open

18  read b
19  c ← a + b
20  write c
```


Thus, this definition is identical to Definition 8, except that the potential conflicting vertex w must occur before v in \mathcal{S} .

The notion of a prefix-race gives rise to an corresponding memory model in which prefix-races are absent.

DEFINITION 11. *The **prefix-race-free** transactional memory model RM_{PFR} is the set of all traces $(C, \sigma) \in \mathcal{U}$ for which there exists a topological sort $\mathcal{S} \in \text{tp} \in$*

beg

of nodes that have finished their execution. The computation tree C and all these associated sets only

We define two time-sensitive sets. The set of *active* transactions at any given time is $\mathbf{active}(C) = \mathbf{ractions}(C)$

7.